

Performance of Bandwidth and QoS Aware LTE Uplink Scheduler Towards Delay Sensitive Traffic

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- ▶ Introduction
- ▶ Scheduler Overview
 - TDPS Metric
 - FDPS Metric
 - Resource Allocation
 - Multi-Bearer User Scheduling
- ▶ Simulation Analysis
 - Single-Bearer Users
 - Double-Bearer Users
 - Triple-Bearer Users
- ▶ Conclusion and Outlook

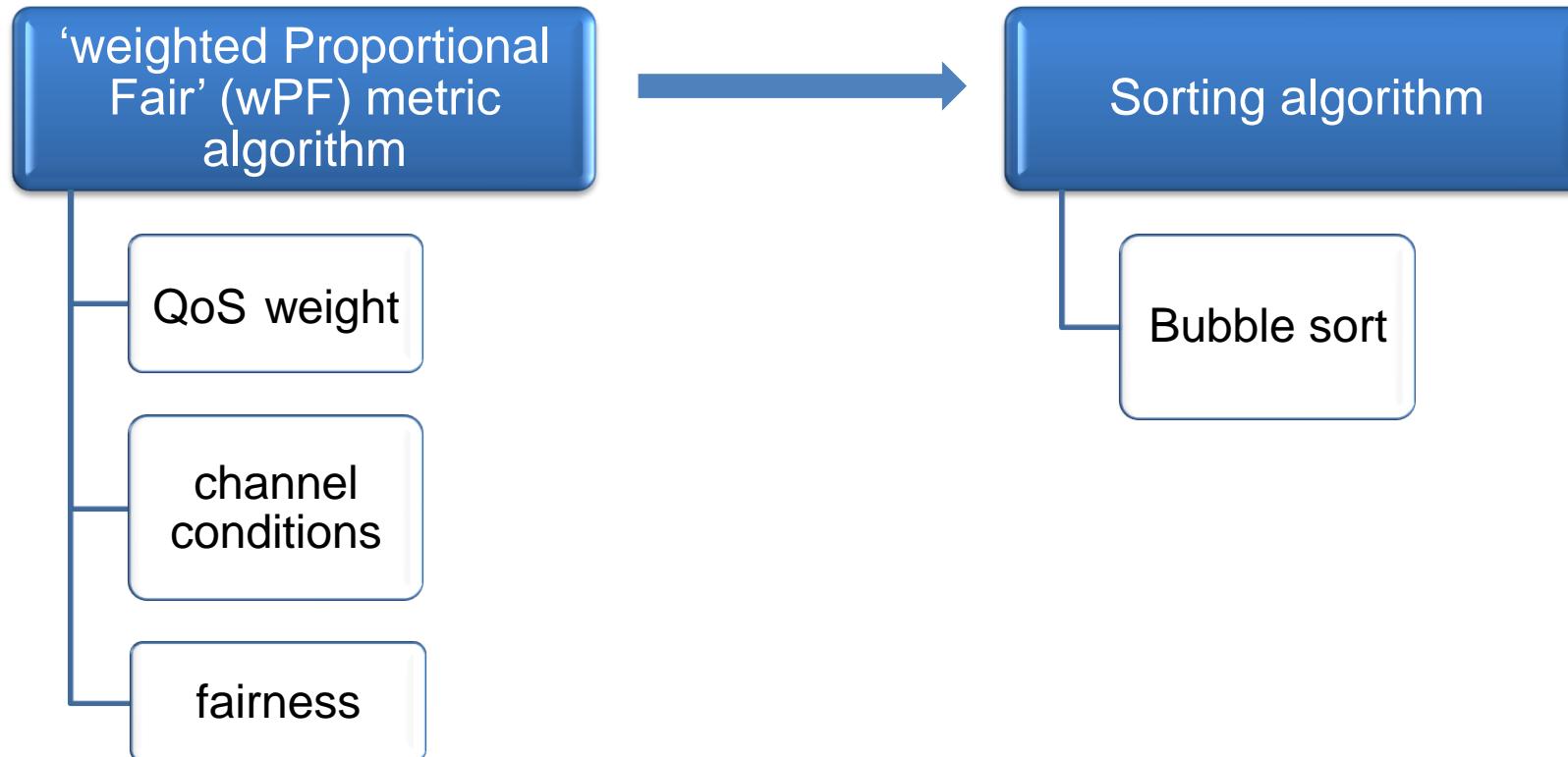
- ▶ Long Term Evolution (LTE) air interface
 - Orthogonal Frequency Division Multiple Access (OFDMA) for downlink
 - Single Carrier Frequency Division Multiple Access (SC-FDMA) for uplink
- ▶ LTE scheduling
 - Allocation of time and frequency resources to users
 - Consider Quality of Service (QoS), channel conditions, fairness, buffer size
- ▶ Uplink scheduling constraints
 - Contiguous Physical Resource Blocks (PRBs) allocation
 - Power Control (PC) based allocation
- ▶ Delay sensitive traffic
 - Stringent delay requirements
 - Scheduling according to QoS demands

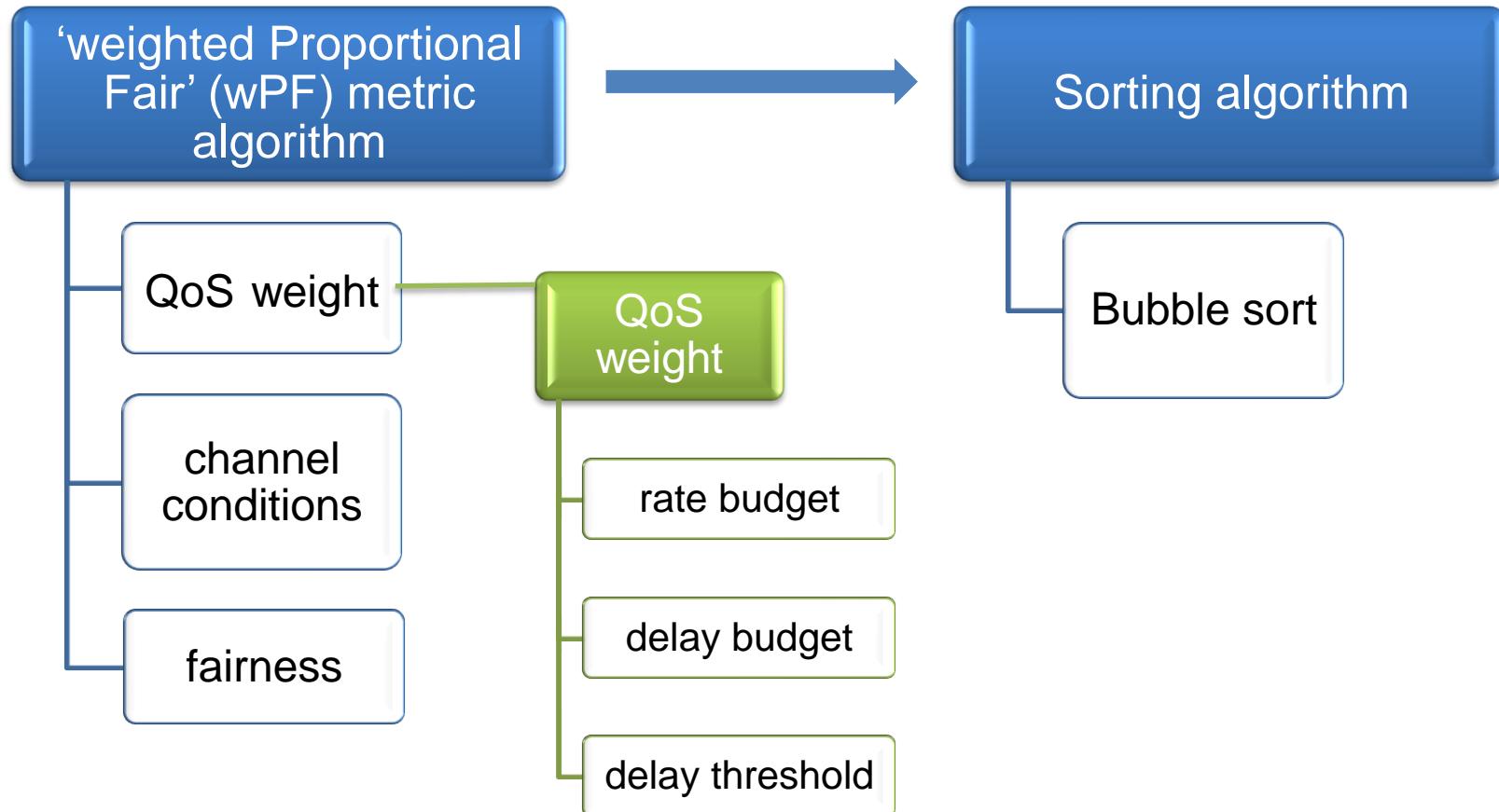
Time Domain Packet Scheduler (TDPS)



Frequency Domain Packet Scheduler (FDPS)







- ▶ All active users in serving eNodeB get TDPS metric values
- ▶ TDPS metric values for user i
 - User QoS weight
 - User instantaneous achievable throughput
 - User Exponential Moving Average (EMA) throughput

$$A_i = \frac{R_{inst,i}(n_i)}{R_{avg,i}} \cdot \sum_k W_{i,k}$$

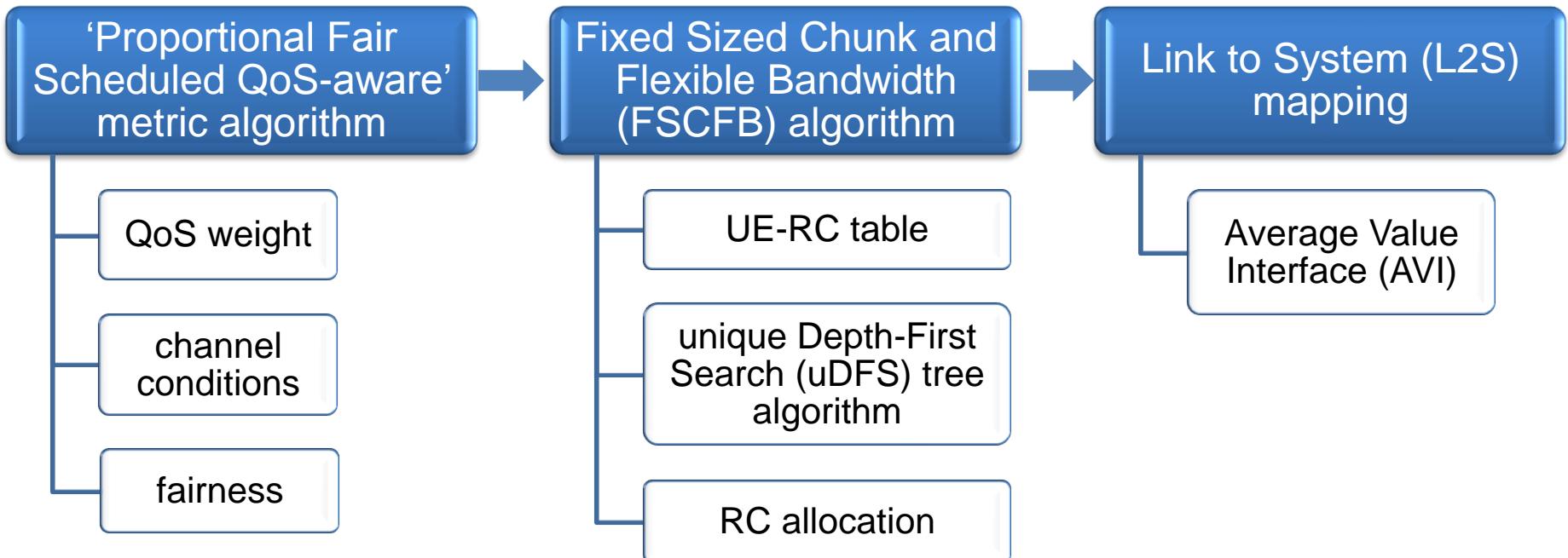
The diagram illustrates the inputs to the TDPS metric calculation. It shows three boxes: 'Channel Conditions', 'QoS Weight', and 'Fairness'. Arrows point from 'Channel Conditions' and 'QoS Weight' to the formula above. An arrow also points from 'Fairness' to the summation term \sum_k .

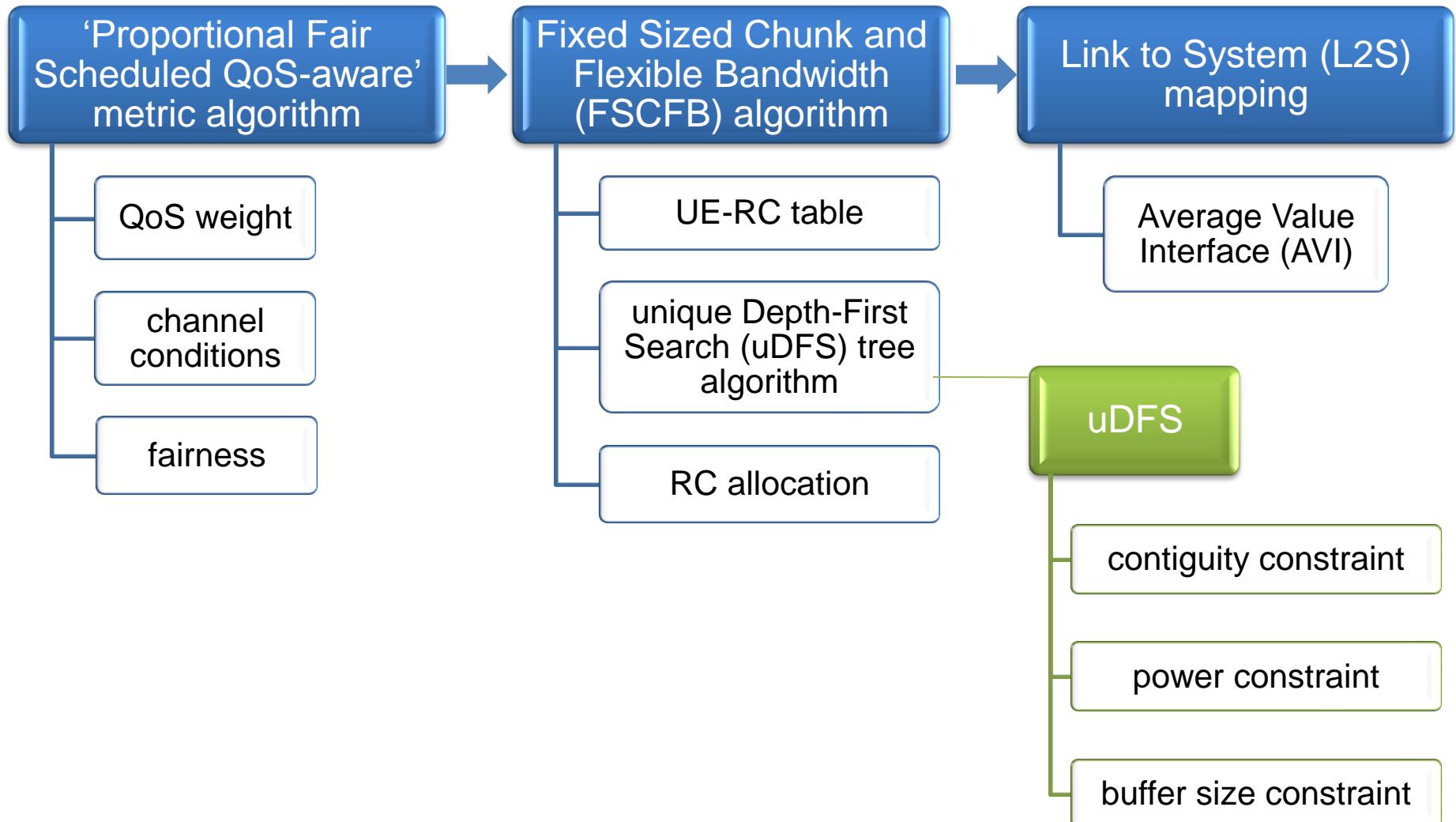
where n_i is maximum number of PRBs user i can get (using PC), $R_{inst,i}$ is the total PRBs in system, $W_{i,k}$ is the QoS weight of bearer k of user i (dynamic weight)

$$W_{i,k} = \frac{R_{min,k}}{R_{avg,i,k}} \cdot \frac{\tau_{i,k}}{\tau_{max,k}} \cdot \rho_k$$

where $R_{min,k}$ and $\tau_{max,k}$ are rate and delay budgets of bearer k , $R_{avg,i,k}$ is the average throughput of k , $\tau_{i,k}$ is the packet delay of k , ρ_k has value 10 if packet delay of k is above threshold, otherwise 1







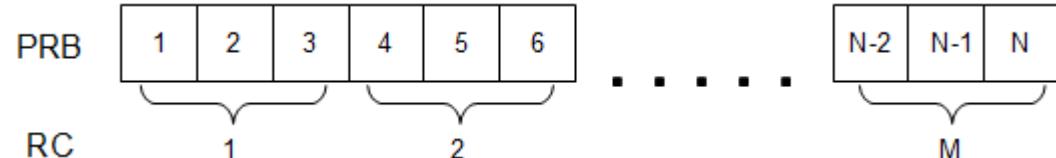
- ▶ All PRBs of all selected users get FDPS metric values
- ▶ FDPS metric value for PRB c of user i
 - User QoS weight
 - PRB instantaneous achievable throughput
 - User EMA throughput over TTIs when user entered FDPS

$$\gamma_i^c = \frac{r_i^c}{R_{sch,avg,i}} \cdot \sum_k W_{i,k}$$

where r_i^c is instantaneously achievable throughput of user i at PRB c and $R_{sch,avg,i}$ is average throughput of user i over the TTIs when it is scheduled by TDPS

► Divide the bandwidth into Resource Chunks (RCs)

- RC is a group of contiguous PRBs
- Algorithm complexity reduced



► Fixed Sized Chunk and Flexible Bandwidth (FSCFB)

- Make UE-RC table of RC metric values
- Determine all possible RC allocation combinations using ‘unique Depth-First Search’ (uDFS) tree algorithm with
 - Contiguity constraint
 - Power constraint
 - Buffer size constraint
- Choose best combination
- Allocate RCs accordingly

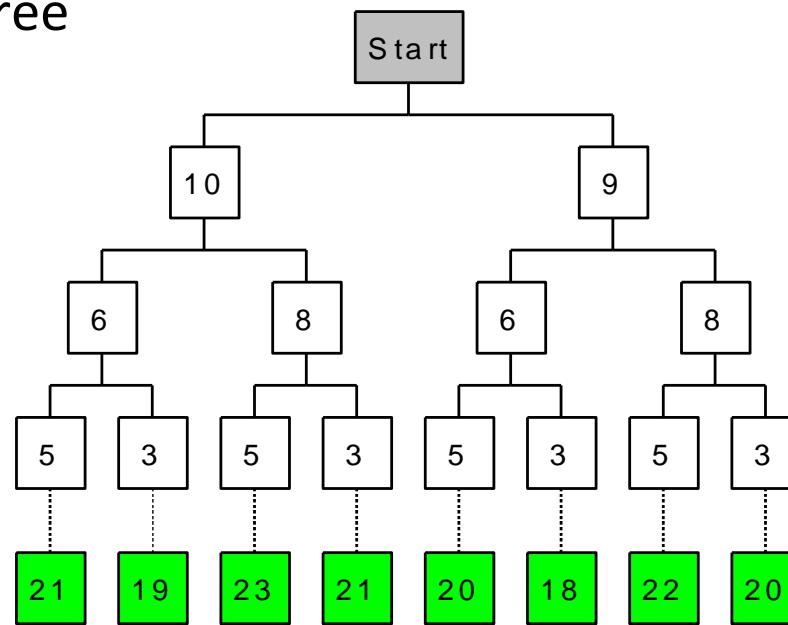
► Average Value Interface (AVI) for user SINR determination

- Take average SINR of the PRBs allocated to a UE
- Determine MCS for UEs according to average SINR
- Determine TBS for UEs according to MCS

- ▶ UE-RC table for 2 users and 3 RCs

	RC_0	RC_1	RC_2
UE_0	10	6	5
UE_1	9	8	3

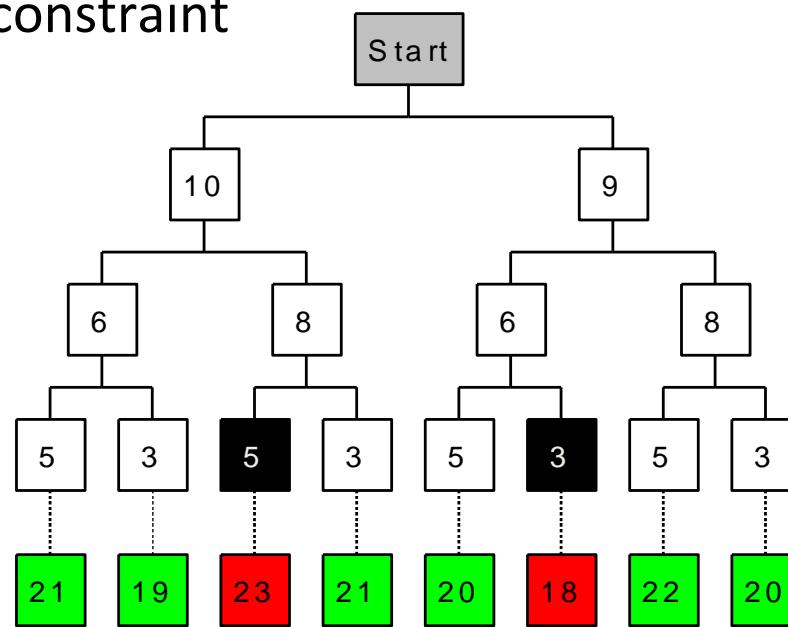
- ▶ Corresponding tree



- ▶ UE-RC table for 2 users and 3 RCs

	RC_0	RC_1	RC_2
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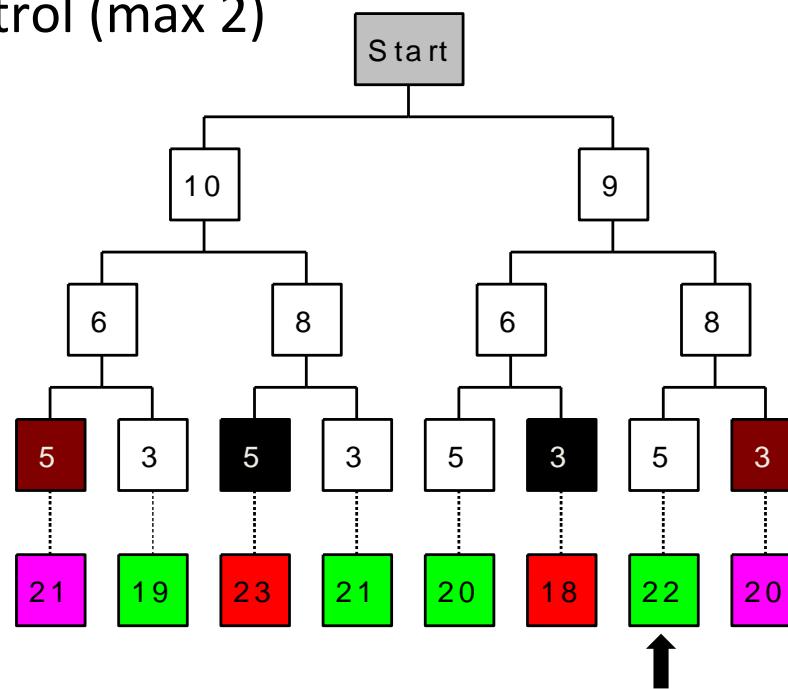
- ▶ With contiguity constraint



- ▶ UE-RC table for 2 users and 3 RCs

	RC_0	RC_1	RC_2
UE_0	10	6	5
UE_1	9	8	3

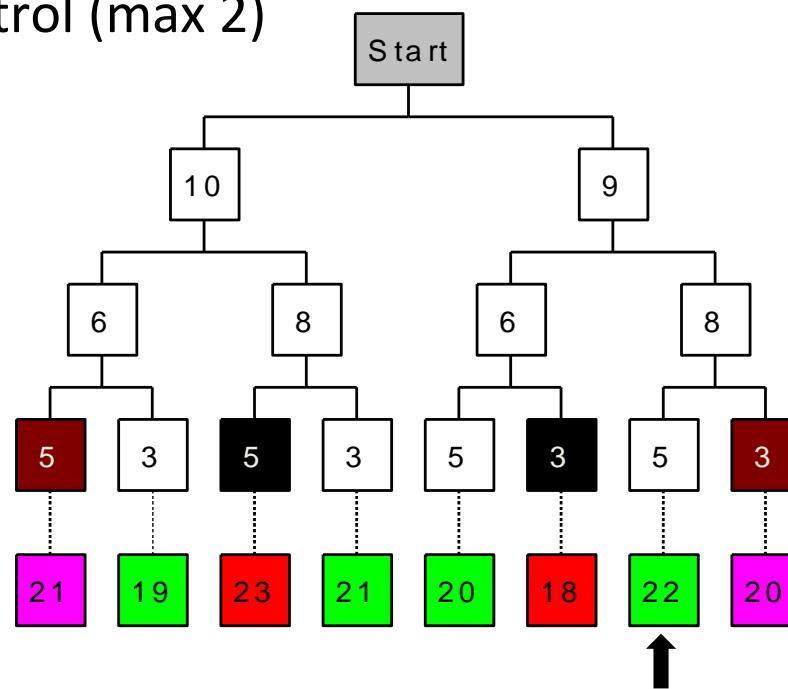
- ▶ With power control (max 2)



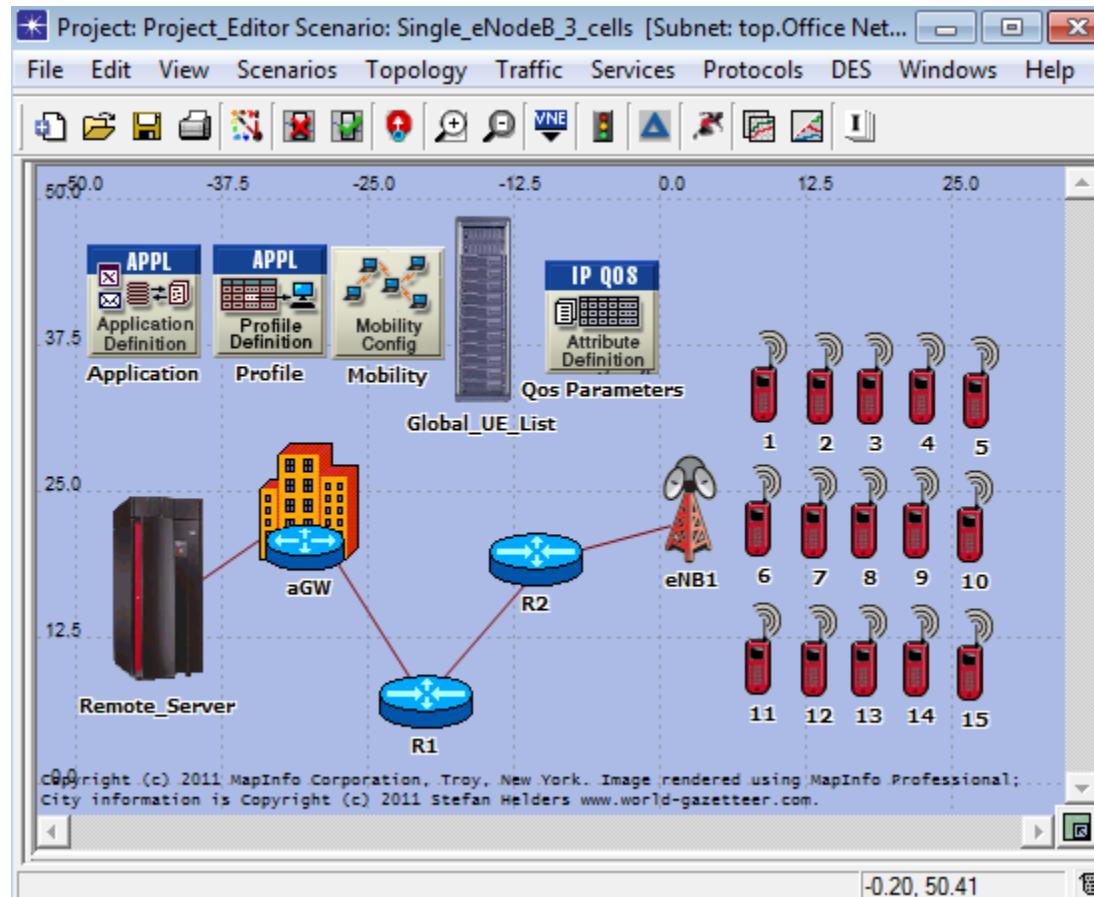
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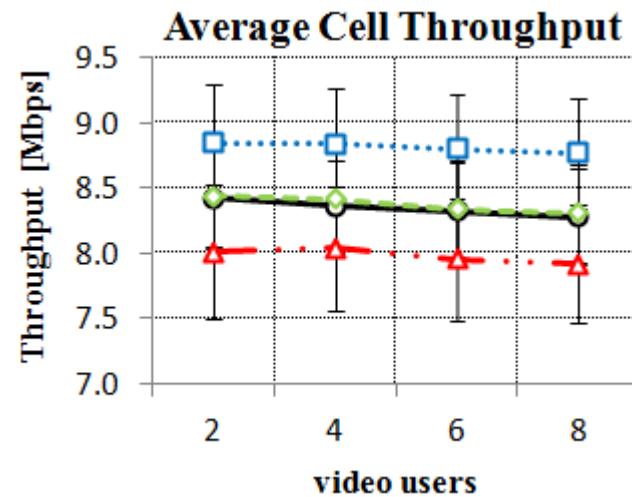
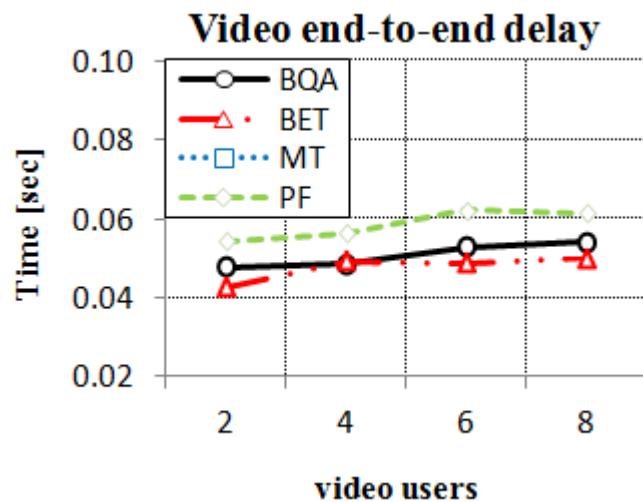
- ▶ PRBs allocation to UE bearers
- ▶ All bearers sorted according to bearer QoS weight
- ▶ PRBs allocated to bearers according to priority
- ▶ Bearers scheduled



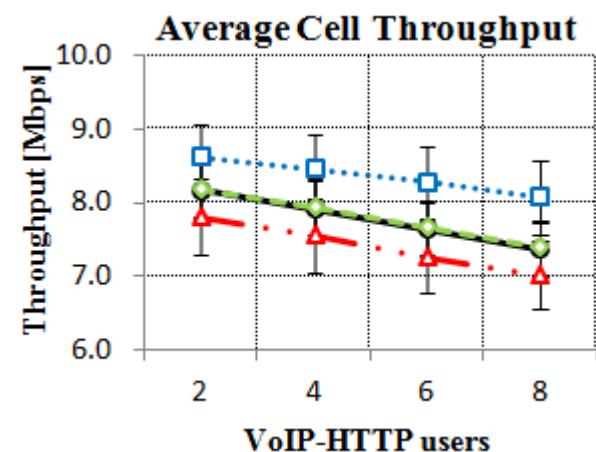
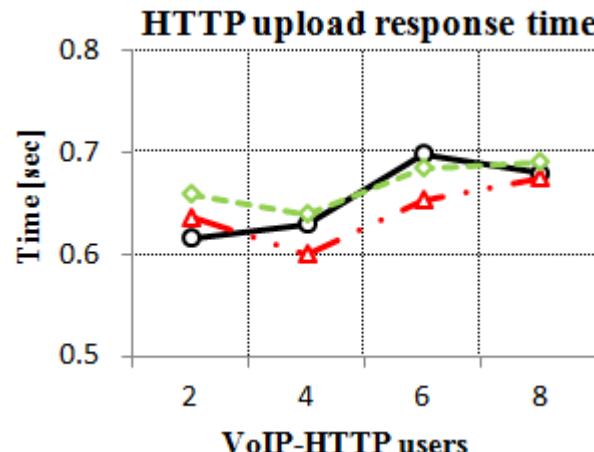
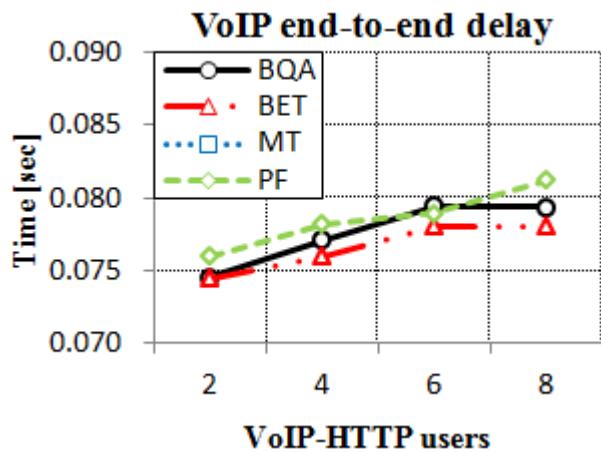
Parameter	Setting
Cell layout	1 eNodeB, 3 cells
System Bandwidth	5 MHz (~25 PRBs)
Frequency reuse factor	1
Cell radius	375m
User velocity	120kmph
Max UE power	23dBm
Path loss	$128.1 + 37.6 \log_{10}(R)$, R in km
Slow fading	Log-normal shadowing, 8dB standard deviation, correlation 1
Fast fading	Jakes-like method
Mobility Model	Random Way Point (RWP)
UE buffer size	Infinite
Power Control	FPC, $\alpha = 0.6$, $P_0 = -58$ dBm
Traffic environment	Loaded
Max FDPS user	5
RC size	5

VoIP traffic model	
Silence/ talk spurt length	Exponential(3) sec
Encoder scheme	GSM EFR
Video traffic model	
Frame size	1200 bytes
Frame inter-arrival time	75ms
HTTP traffic model	
Page size	100Kbytes
Page inter-arrival time	12 sec
FTP traffic model	
File size	20Mbytes
File inter-request time	Uniform distribution, min 80s, max 100s

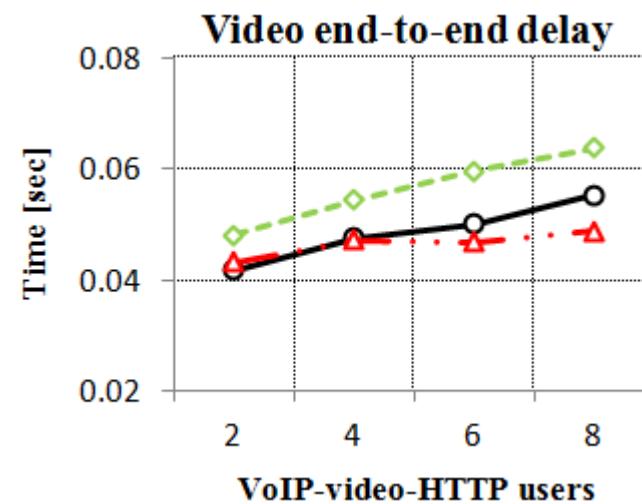
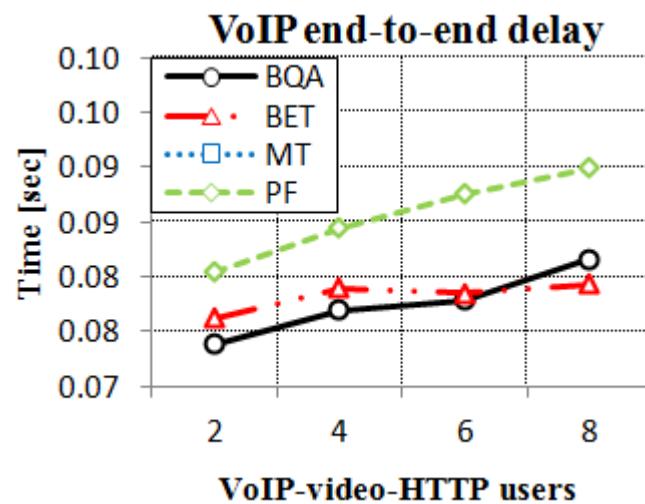
- ▶ Fixed 8 FTP users to keep the cell loaded
- ▶ Single-bearer video users added stepwise into the loaded cell



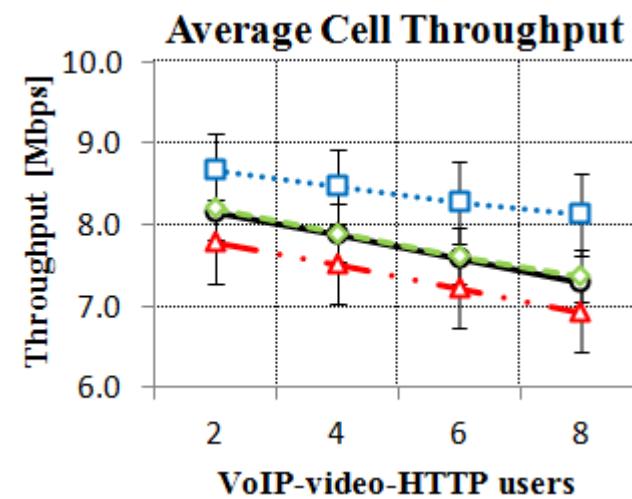
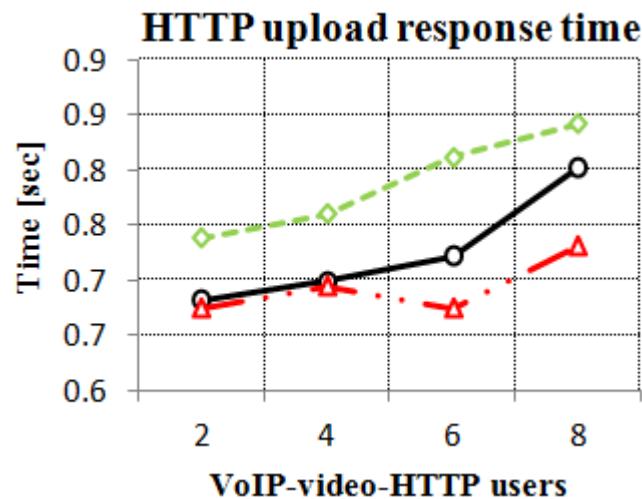
- ▶ Fixed 8 FTP users to keep the cell loaded
- ▶ Double-bearer VoIP-HTTP users added stepwise into the loaded cell



- ▶ Fixed 8 FTP users to keep the cell loaded
- ▶ Triple-bearer VoIP-video-HTTP users added stepwise into the loaded cell



- ▶ Fixed 8 FTP users to keep the cell loaded
- ▶ Triple-bearer VoIP-video-HTTP users added stepwise into the loaded cell



► BQA scheduler evaluation

- Delay sensitive traffic QoS guarantee
- Comparison with contemporary schedulers
- All schedulers utilize
 - Channel conditions
 - FSCFB (and uDFS) algorithm
 - Multi-bearer user scheduling
- Dynamic QoS weight is the distinctive feature of BQA

► Future goals

- Integrate HARQ functionality
- Include Admission Control (AC) functionality

THANK YOU!

Questions???